

# Subchannel Allocation for Vehicle-to-Vehicle Broadcast Communications in Mode-3

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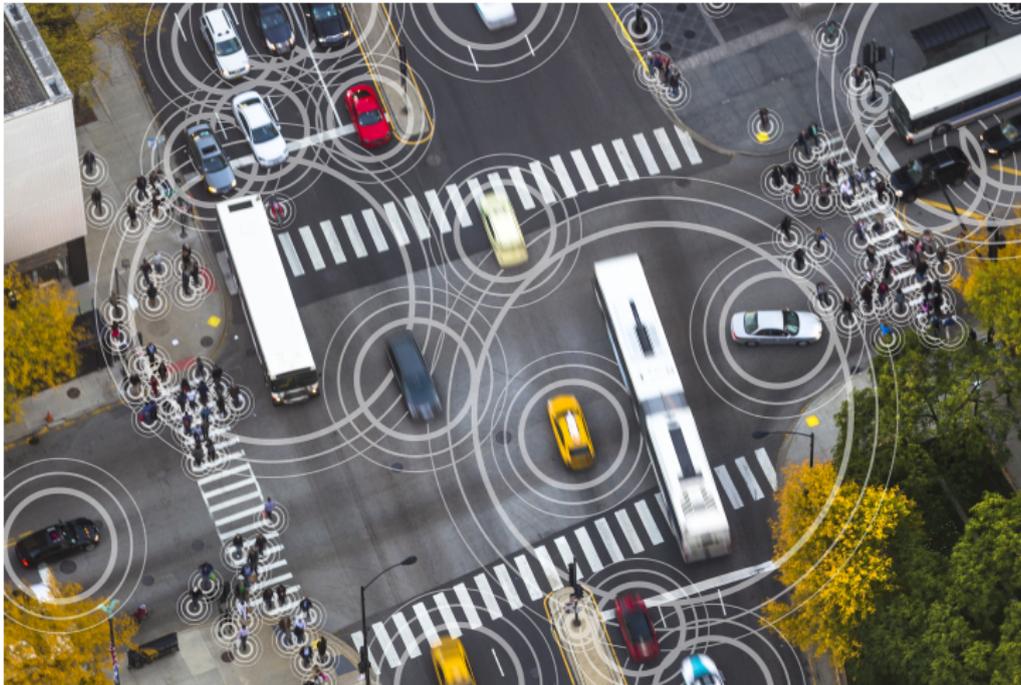


Figure 1: Connected world

# Background

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- 3GPP<sup>1</sup> proposed in Release 14, two novel schemes to support sidelink vehicular communications
  - C-V2X *mode-3* (centralized)
  - C-V2X<sup>2</sup> *mode-4* (distributed)

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<sup>1</sup>The 3rd Generation Partnership Project

<sup>2</sup>Cellular Vehicle-to-Everything

<sup>3</sup>Device-to-Device (D2D) communications

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  - C-V2X *mode-3* (centralized)
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- C-V2X *modes* are based on LTE-D2D<sup>3</sup> technology, where similar communication modalities were proposed.

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  - C-V2X *mode-3* (centralized)
  - C-V2X<sup>2</sup> *mode-4* (distributed)
- C-V2X *modes* are based on LTE-D2D<sup>3</sup> technology, where similar communication modalities were proposed.
- However, in LTE-D2D (initially introduced for PS) the most important criterion was to prolong batteries lifespan (at the expense of compromising latency).

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- To fulfill the low latency and high reliability requirements:

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<sup>5</sup>A subchannel is a time-frequency resource chunk.

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- Modifications at PHY layer
  - Denser distribution of DMRS<sup>4</sup>

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# Background

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- To fulfill the low latency and high reliability requirements:
- Modifications at PHY layer
  - Denser distribution of DMRS<sup>4</sup>
- Modifications at MAC layer
  - A novel subchannelization<sup>5</sup> containing
    - (i) scheduling assignments (SCI)
    - (ii) data (TB)in the same subframe to minimize latency.

A semi-persistent scheduling was proposed for mode-4. No approach has been specified for mode-3.

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<sup>5</sup>A subchannel is a time-frequency resource chunk.

# Background

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- Besides **uplink** and **downlink** (Uu), vehicles can also communicate via **sidelink** (PC5), which supports direct communications between vehicles.

# C-V2X Mode-3

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- Conversely to mainstream communications, in **C-V2X mode-3** data traffic from/to vehicles do not traverse the eNodeB.

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  - eNodeBs **only** intervene in the **resource allocation** process.

# C-V2X Mode-3

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  - eNodeBs **only** intervene in the **resource allocation** process.
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# C-V2X Mode-3

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- Conversely to mainstream communications, in **C-V2X mode-3** data traffic from/to vehicles do not traverse the eNodeB.
  - eNodeBs **only** intervene in the **resource allocation** process.
  - Then **vehicles communicate directly** with their counterparts via sidelink
- In **safety** applications, vehicles would typically exchange **cooperative awareness messages (CAMs)**: position, velocity, direction, etc.

# C-V2X Mode-3

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# C-V2X Mode-3

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- As these messages transport important information, it is crucial that they are received **reliably**.
- Due to the one-to-all broadcast nature, the allocation of resources (or subchannels) slightly differs from mainstream communications.
- **Example:** *If two vehicles transmit concurrently they will not receive the CAM message of the other.*
- *Four types of conflicts/requirements have been identified.*

# Problem Identification: Condition Type I

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## Condition Type I: Differentiated QoS Requirements per Vehicle

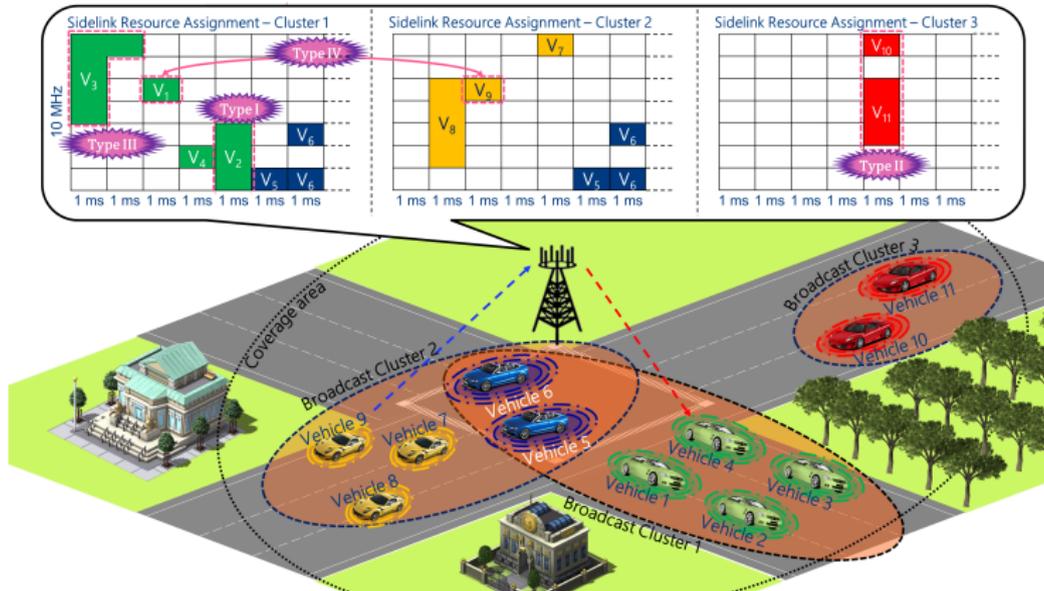


Figure 2: Different types of allocation conflicts

# Problem Identification: Condition Type II

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## Condition Type II: Intra-cluster Subframe Allocation Conflicts

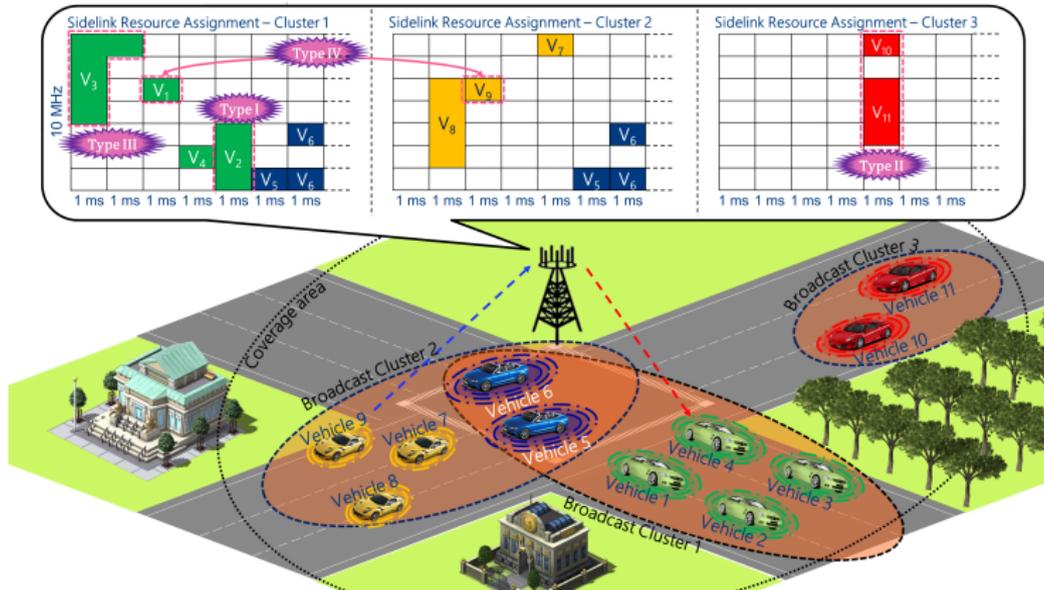


Figure 3: Different types of allocation conflicts

# Problem Identification: Condition Type III

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## Condition Type III: Minimal Time Dispersion of Subchannels

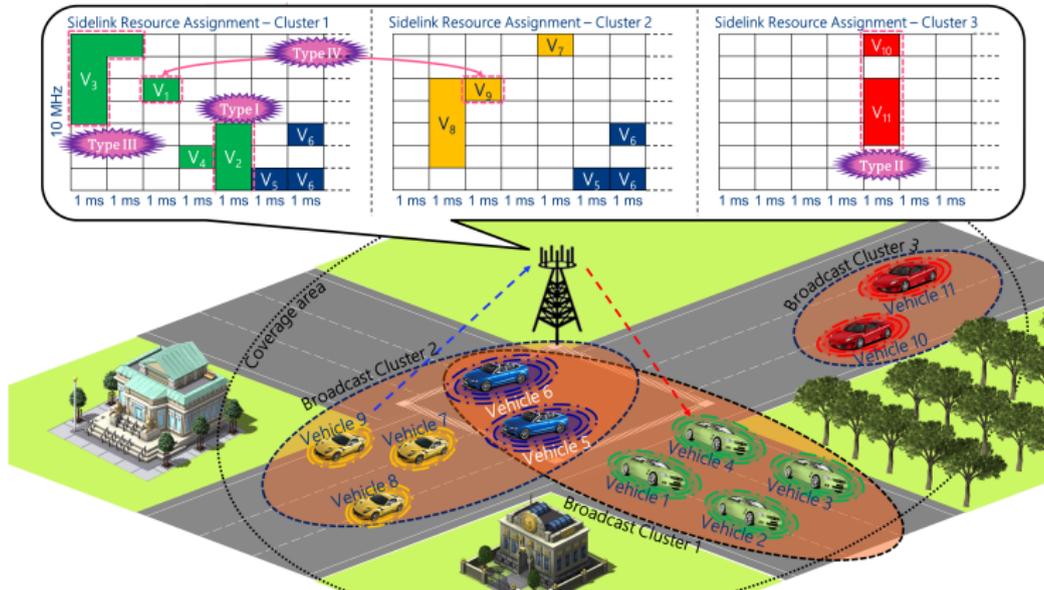


Figure 4: Different types of allocation conflicts

# Problem Identification: Condition Type IV

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## Condition Type IV: One-hop Inter-cluster Subchannel Conflicts

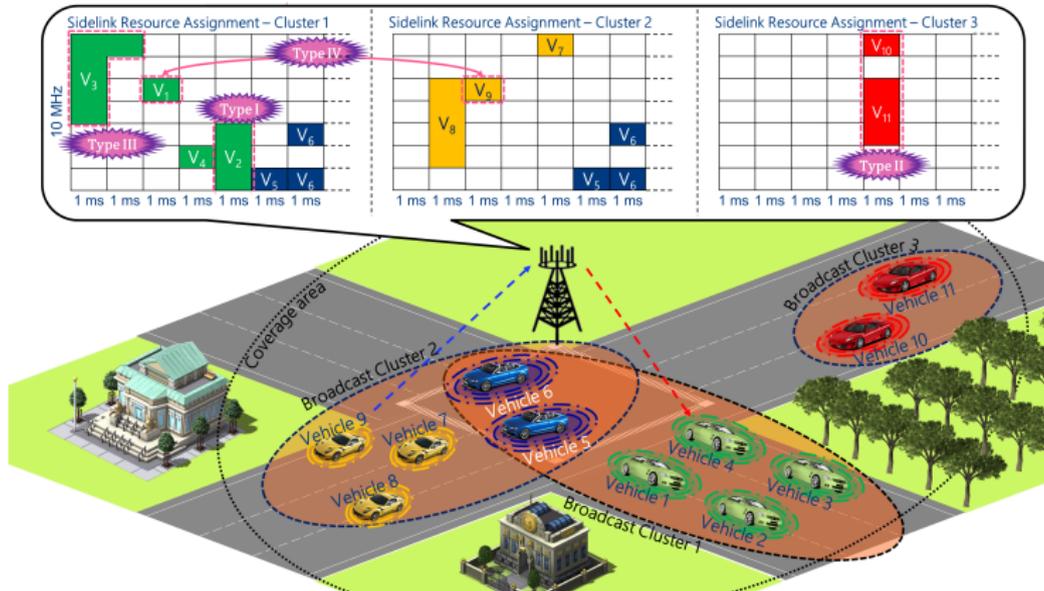
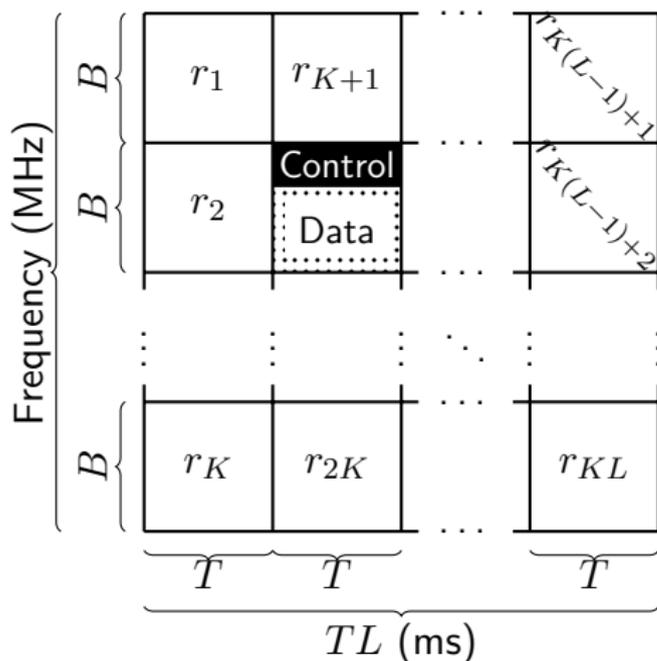


Figure 5: Different types of allocation conflicts

## Sidelink Channelization

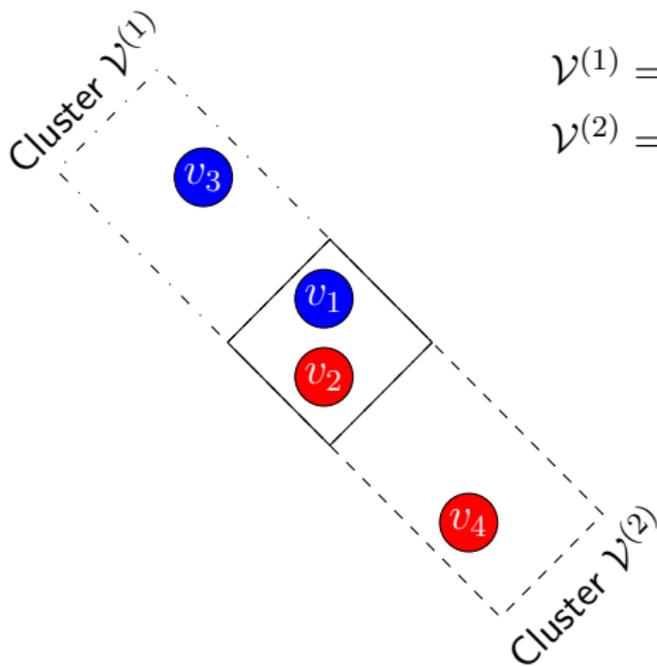
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- $T$ : duration of a subframe
- $K$ : number of subchannels per subframe
- $L$ : total number of subframes for allocation
- $B$ : subchannel bandwidth

## Motivation: Toy Example

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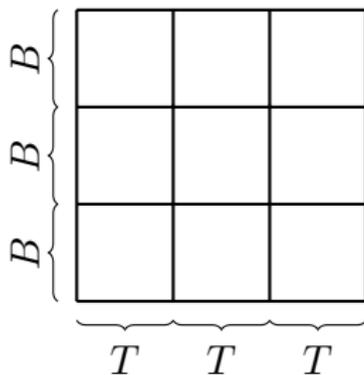
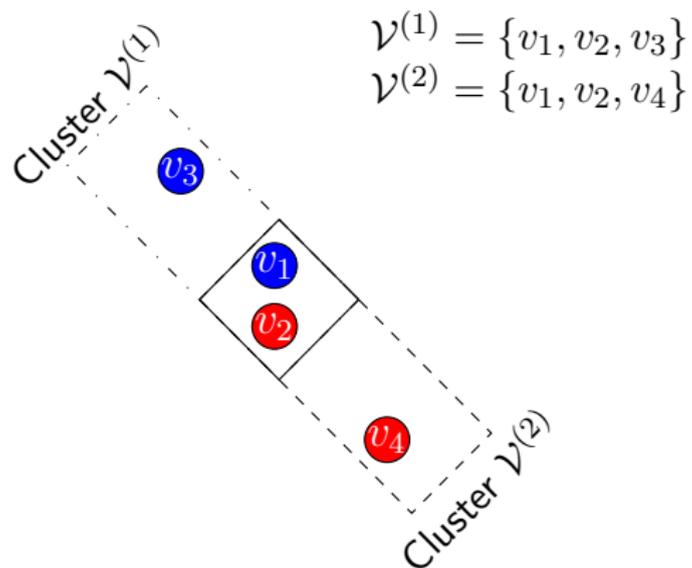


$$\mathcal{V}^{(1)} = \{v_1, v_2, v_3\}$$

$$\mathcal{V}^{(2)} = \{v_1, v_2, v_4\}$$

## Motivation: Toy Example

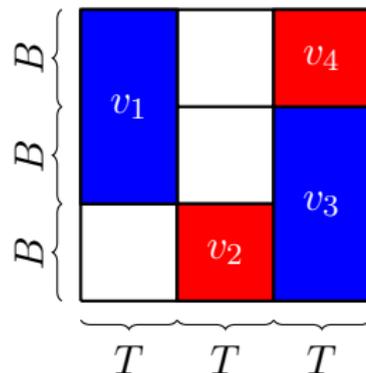
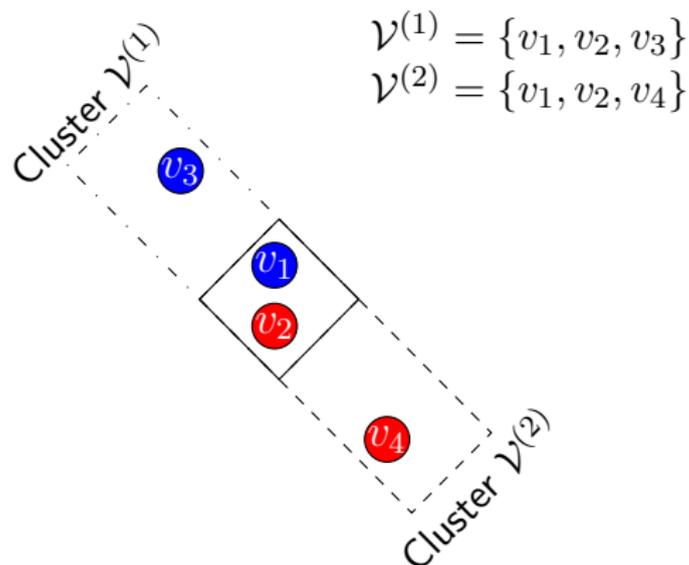
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# Motivation: Toy Example

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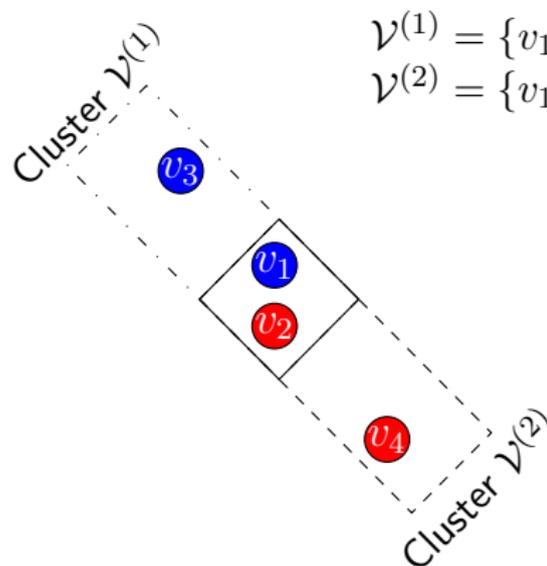
Condition Type I: Differentiated QoS Requirements per Vehicle



## Motivation: Toy Example

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Condition Type II: Intra-cluster Subframe Allocation Conflicts



$$\mathcal{V}^{(1)} = \{v_1, v_2, v_3\}$$

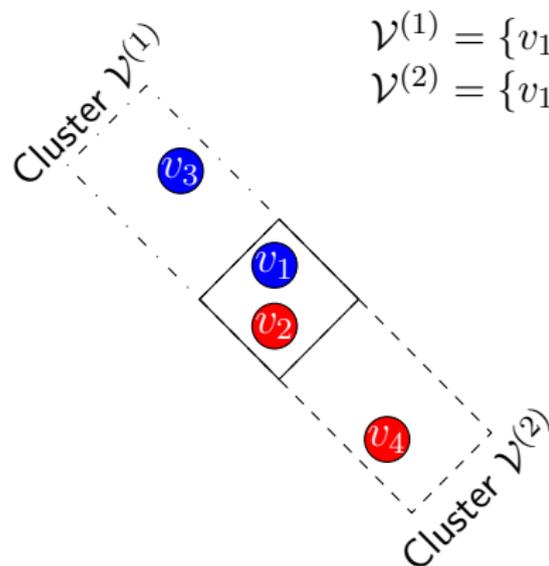
$$\mathcal{V}^{(2)} = \{v_1, v_2, v_4\}$$

$$\tilde{\mathbf{G}} = \begin{bmatrix} v_1 & v_2 & v_3 & v_4 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{matrix}$$

## Motivation: Toy Example

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Condition Type III: Minimal Time Dispersion of Subchannels



$$\mathcal{V}^{(1)} = \{v_1, v_2, v_3\}$$

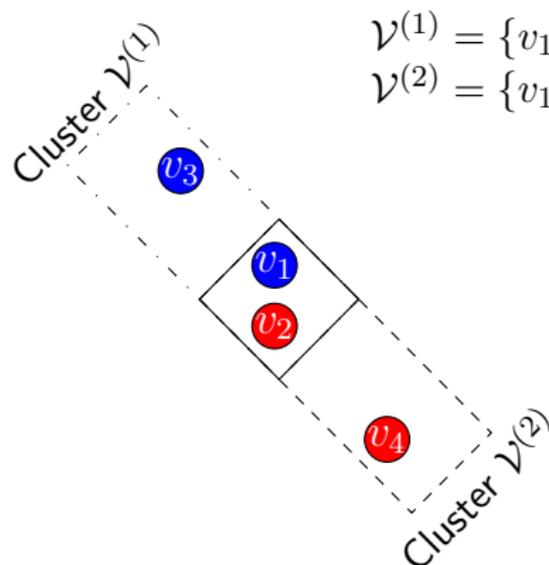
$$\mathcal{V}^{(2)} = \{v_1, v_2, v_4\}$$

$$\tilde{\mathbf{Q}} = \begin{bmatrix} sf_1 & sf_2 & sf_3 \\ 0 & 0 & 0 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \end{bmatrix} \begin{matrix} sf_1 \\ sf_2 \\ sf_3 \end{matrix}$$

## Motivation: Toy Example

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Condition Type IV: One-hop Inter-cluster Subchannel Conflicts



$$\mathcal{V}^{(1)} = \{v_1, v_2, v_3\}$$

$$\mathcal{V}^{(2)} = \{v_1, v_2, v_4\}$$

$$\tilde{\mathbf{H}} = \begin{bmatrix} v_1 & v_2 & v_3 & v_4 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix} \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{matrix}$$

# Problem Formulation

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$$\max \mathbf{c}^T \mathbf{x}$$

subject to

$$\mathbf{q}_{N \times 1} - \epsilon \leq (\mathbf{I}_{N \times N} \otimes \mathbf{1}_{1 \times KL})(\mathbf{c}_{NKL \times 1} \circ \mathbf{x}_{NKL \times 1}) \leq \mathbf{q}_{N \times 1} + \epsilon$$

$$[(\mathbf{G}_{P \times N}^+ \otimes \mathbf{I}_{L \times L})(\mathbf{I}_{NL \times NL} \otimes \mathbf{1}_{1 \times K})\mathbf{x}] \circ [(\mathbf{G}_{P \times N}^- \otimes \mathbf{I}_{L \times L})(\mathbf{I}_{NL \times NL} \otimes \mathbf{1}_{1 \times K})\mathbf{x}] = \mathbf{0}_{PL \times 1}$$

$$[(\mathbf{I}_{N \times N} \otimes \mathbf{Q}_{L \times L}^+)(\mathbf{I}_{NL \times NL} \otimes \mathbf{1}_{1 \times K})\mathbf{x}] \circ [(\mathbf{I}_{N \times N} \otimes \mathbf{Q}_{L \times L}^-)(\mathbf{I}_{NL \times NL} \otimes \mathbf{1}_{1 \times K})\mathbf{x}] = \mathbf{0}_{NL \times 1}$$

$$[(\mathbf{H}_{U \times N}^+ \otimes \mathbf{I}_{KL \times KL})\mathbf{x}] \circ [(\mathbf{H}_{U \times N}^- \otimes \mathbf{I}_{KL \times KL})\mathbf{x}] = \mathbf{0}_{U \times 1}.$$

# Relaxed Formulation

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$$\max \mathbf{c}^T \mathbf{x}$$

subject to

$$\mathbf{q}_{N \times 1} - \epsilon \leq (\mathbf{I}_{N \times N} \otimes \mathbf{1}_{1 \times KL})(\mathbf{c}_{NKL \times 1} \circ \mathbf{x}_{NKL \times 1}) \leq \mathbf{q}_{N \times 1} + \epsilon$$

$$\mathbf{x}^T (\mathbf{I}_{NL \times NL} \otimes \mathbf{1}_{K \times 1}) \{ \tilde{\mathbf{G}}_{N \times N} \otimes \mathbf{I}_{L \times L} \} (\mathbf{I}_{NL \times NL} \otimes \mathbf{1}_{1 \times K}) \mathbf{x} = 0$$

$$\mathbf{x}^T (\mathbf{I}_{NL \times NL} \otimes \mathbf{1}_{K \times 1}) \{ \mathbf{I}_{N \times N} \otimes \tilde{\mathbf{Q}}_{L \times L} \} (\mathbf{I}_{NL \times NL} \otimes \mathbf{1}_{1 \times K}) \mathbf{x} = 0$$

$$\mathbf{x}^T \{ \tilde{\mathbf{H}}_{N \times N} \otimes \mathbf{I}_{KL \times KL} \} \mathbf{x} = 0.$$

# Simulation Scenario

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Consider the following setting:

- There is a total of  $N = 40$  vehicles divided into 4 clusters:  
 $|\mathcal{V}^{(1)}| = 16$     $|\mathcal{V}^{(2)}| = 16$     $|\mathcal{V}^{(3)}| = 16$     $|\mathcal{V}^{(4)}| = 8$   
such that:
  - $|\mathcal{V}^{(1)} \cap \mathcal{V}^{(2)} \cap \mathcal{V}^{(3)}| = 8$
  - $|\mathcal{V}^{(1)} \cap \mathcal{V}^{(4)}| = \emptyset$
  - $|\mathcal{V}^{(2)} \cap \mathcal{V}^{(4)}| = \emptyset$
  - $|\mathcal{V}^{(3)} \cap \mathcal{V}^{(4)}| = \emptyset$
- QoS requirements: 12 Mbps, 10 Mbps, 5 Mbps or 3 Mbps.
- There are 10 vehicles for each kind of QoS.

# Scenario : Required QoS = 12 Mbps

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The number of subframes is  $L = 16$ .

The number of subchannels per subframe is  $K = 4$ .

$\epsilon = 0.8$  Mbps and therefore the range of rates are

[11.2 – 12.8] Mbps, [9.2 – 10.8] Mbps, [4.2 – 5.8] Mbps and [2.2 – 3.8] Mbps.

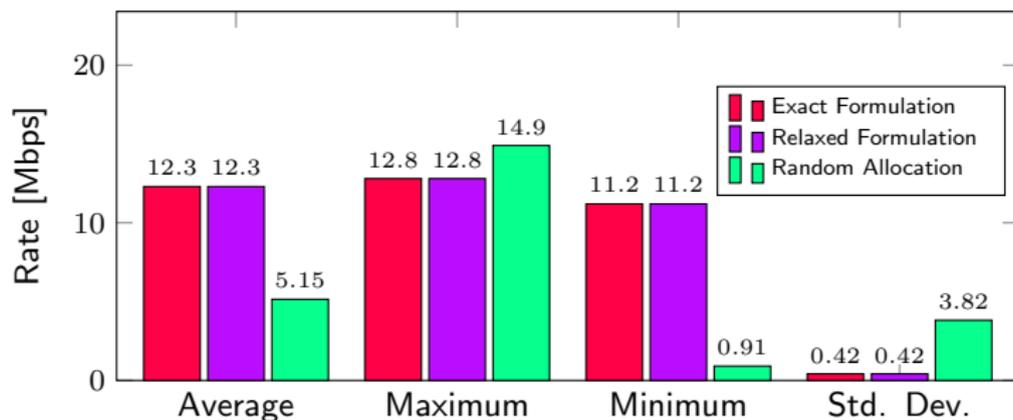


Figure 6: Scenario 1 / Vehicles with QoS = 12 Mbps

# Scenario: Required QoS = 10 Mbps

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The number of subframes is  $L = 16$ .

The number of subchannels per subframe is  $K = 4$ .

$\epsilon = 0.8$  Mbps and therefore the range of rates are

$[11.2 - 12.8]$  Mbps,  $[9.2 - 10.8]$  Mbps,  $[4.2 - 5.8]$  Mbps and  $[2.2 - 3.8]$  Mbps.

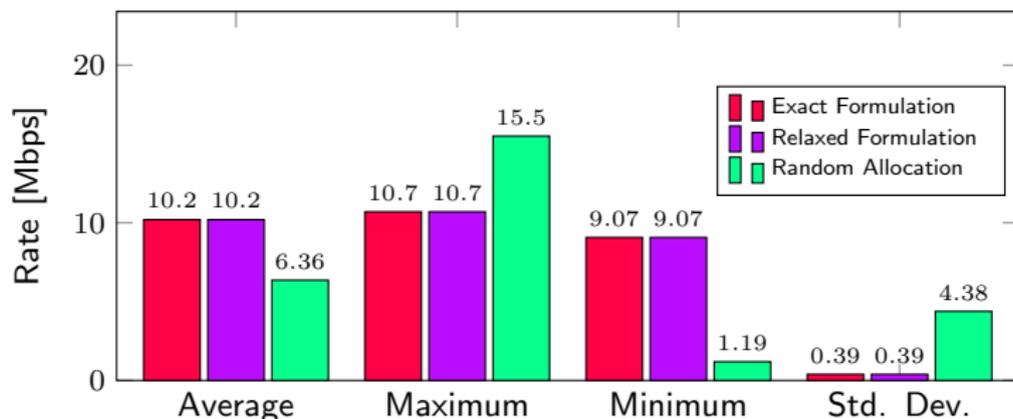


Figure 7: Scenario 1 / Vehicles with QoS = 10 Mbps

# Scenario: Required QoS = 5 Mbps

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The number of subframes is  $L = 16$ .

The number of subchannels per subframe is  $K = 4$ .

$\epsilon = 0.8$  Mbps and therefore the range of rates are

$[11.2 - 12.8]$  Mbps,  $[9.2 - 10.8]$  Mbps,  $[4.2 - 5.8]$  Mbps and  $[2.2 - 3.8]$  Mbps.

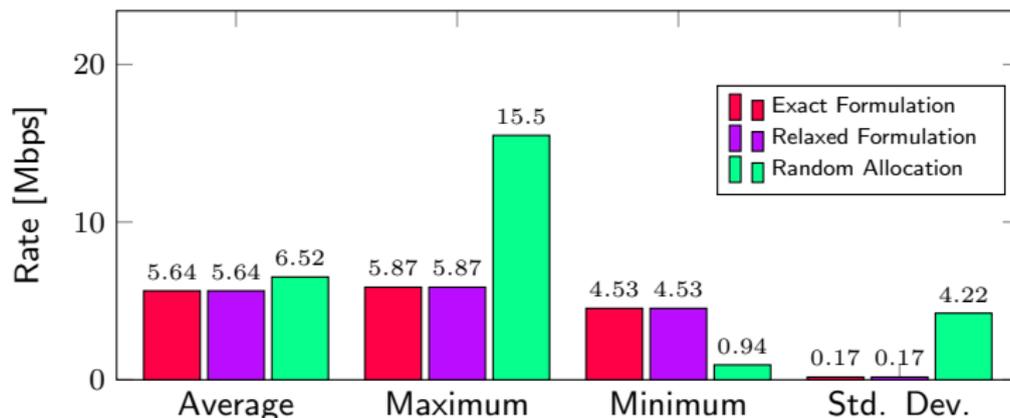


Figure 8: Scenario 1 / Vehicles with QoS = 5 Mbps

# Scenario: Required QoS = 3 Mbps

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The number of subframes is  $L = 16$ .

The number of subchannels per subframe is  $K = 4$ .

$\epsilon = 0.8$  Mbps and therefore the range of rates are

$[11.2 - 12.8]$  Mbps,  $[9.2 - 10.8]$  Mbps,  $[4.2 - 5.8]$  Mbps and  $[2.2 - 3.8]$  Mbps.

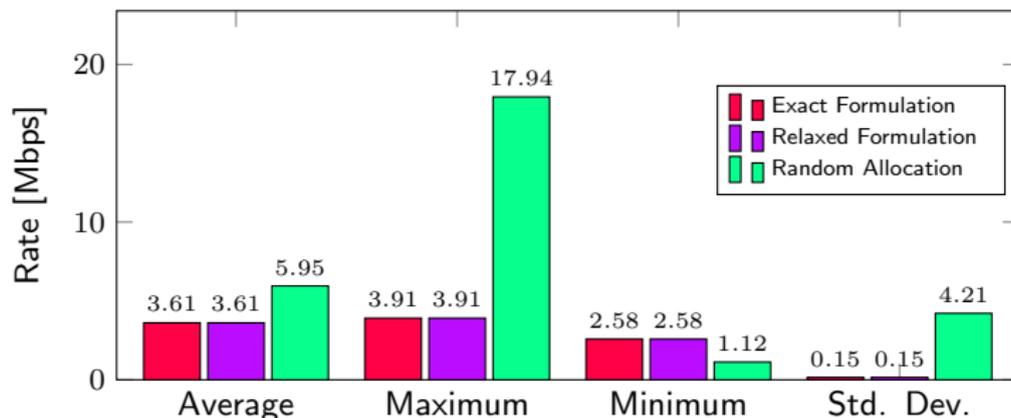


Figure 9: Scenario 1 / Vehicles with QoS = 3 Mbps

# Conclusions

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- In this work we have presented a subchannel allocation framework for C-V2X *mode-3*.
- Four types of conditions have been identified and incorporated in order to guarantee a conflict-free allocation that complies with QoS requirements per vehicle.
- In addition, a relaxed formulation (RF) of the original problem that does not impinge on optimality was proposed
- The QoS requirements can be very tightly met—with the exact formulation and the relaxed version—but the random approach introduces noticeable deviation.

# Questions

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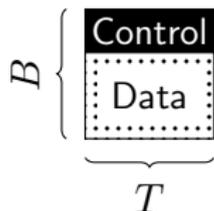
Email: [l.f.abanto@tue.nl](mailto:l.f.abanto@tue.nl)

# Subchannel Structure

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Assuming a 10 MHz ITS (Intelligent Transportation Systems) channel, up to 7 subchannels per subframe can be obtained. Thus,

- $B$ : 1.26 MHz
- $T$ : 1 ms (2 slots of 0.5 ms each)
- Control: 2 RBs<sup>6</sup> per slot  $\leftarrow$  24 subcarriers
- Data: 5 RBs per slot  $\leftarrow$  60 subcarriers



## Subchannel

A subchannel of 7 RBs is capable of transporting a basic CAM message with a payload of 200 bytes.

<sup>6</sup>RB: A resource block consists of 12 subcarriers

# Properties

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## Property 1 (Product of two tensor products)

Let  $\mathbf{X} \in \mathbb{R}^{m \times n}$ ,  $\mathbf{Y} \in \mathbb{R}^{r \times s}$ ,  $\mathbf{W} \in \mathbb{R}^{n \times p}$ , and  $\mathbf{Z} \in \mathbb{R}^{s \times t}$ , then

$$\mathbf{X}\mathbf{Y} \otimes \mathbf{W}\mathbf{Z} = (\mathbf{X} \otimes \mathbf{W})(\mathbf{Y} \otimes \mathbf{Z}) \in \mathbb{R}^{mr \times pt}$$

## Property 2 (Pseudo-inverse of a tensor product)

Let  $\mathbf{X} \in \mathbb{R}^{m \times n}$  and  $\mathbf{Y} \in \mathbb{R}^{r \times s}$ , then

$$(\mathbf{X} \otimes \mathbf{Y})^\dagger = \mathbf{X}^\dagger \otimes \mathbf{Y}^\dagger \in \mathbb{R}^{ns \times mr}$$